# Application of F-G Diagonalization Algorithm to Restricted Maximum Likelihood Estimation of Variance Components 

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#### Abstract

The F-G algorithm of Flury and Gautschi can be used to find an orthogonal matrix B such that:  positive weights. The orthogonal matrix $B$ can be interpreted as the matrix which brings matrices $C_{1} \ldots \ldots$, $\mathrm{C}_{\mathrm{k}}$ simultaneously as close to diagonality as possible. To reduce the number of operations required by $\mathrm{F}-\mathrm{G}$ algorithm, Clarkson used a modified algorithm (MF-G) to find an orthogonal matrix B such that $\mathrm{B}^{\prime} \mathrm{C}_{1} \mathrm{~B}$ is nearly diagonal. Both F-G and MF-G algorithm were applied to three sets of mixed model coefficient matrices in animal breeding cases. Close estimate to the exact REML solutions were obtained for traits with low heritability (large $x$ ). One can use equal or unequal weights $n_{1}, \ldots \ldots, n_{k}$ to achieve convergence for both algorithms.


## Introduction

Variance component estimation can be very demanding computationally for large data set. Restricted maximum likelihood (REML) was derived by Patterson and Thompson [1] whose purpose was to eliminate the bias in maximum likelihood (ML) due to estimation of fixed effects. Smith and Graser [2] described an efficient algorithm for computing REML estimators of variance components in a class of mixed model. They tridiagonalized the coefficient matrix through a series of Householder transformations so that direct inversion of the coefficient matrix was unnecessary. They found that evaluation of $\operatorname{tr}\left(\mathrm{Z}^{\prime} \mathrm{MZ}+\propto \mathrm{I}\right)^{-1}$ is the most computationaly demanding step in the tridiagonalization procedure. However, computing this trace becomes a computational triviality using procedure through singular value decom-
position. Since ( $\left.Z^{\prime} M Z+\infty \mathrm{I}\right)$ and $(\mathrm{D}+\infty \mathrm{I})$ are similar matrices where $\mathrm{D}=\mathrm{V}^{\prime}\left(\mathrm{Z}^{\prime} \mathrm{MZ}\right) \mathrm{V}$ and V is an orthogonal matrix and, hence, $\operatorname{tr}\left(\mathrm{Z}^{\prime} \mathrm{MZ}+\propto \mathrm{I}\right)^{-1}=$ $(\mathrm{D}+\propto \mathrm{I})^{-1}$. So $\operatorname{Tr}\left(\mathrm{Z}^{\prime} M Z+\infty \mathrm{I}\right)^{-1}$ is simply the sum of the reciprocals of the diagonal elements of the matrix ( $\mathrm{D}+\infty \mathrm{I}$ ).

In animal breeding, if the sires are related with a relationship matrix (A), Smith and Graser [2] suggested to redefine $Z_{1}$ as $Z L$ such that $A=L L^{\prime}$ where $L$ is the lower triangular matrix obtained by applying Cholesky Decomposition to A. So $\left(Z^{\prime}{ }_{1} M Z+\infty I\right) s^{*}=Z_{1} M Y$ where $M=I-X^{\prime}\left(X^{\prime} X\right)^{-1} X$ and $s^{*}=L^{-1} s$.

Patterson and Thompson [1] and Thompson and Cameron [3] suggested the diagonalization of the coefficient matrix ( $Z^{\prime} M Z+\infty I$ ) to reduce the CPU time required to obtain direct inverse in each iteration. Their basic idea was to calculate the inverse of $\left(Z^{\prime} M Z+\propto I\right)$ by computing $V(D+\propto I)^{-1} V^{\prime}$ instead of direct inversion because $\left(\mathrm{Z}^{\prime} \mathrm{MZ}+\propto \mathrm{I}\right)^{-1}=\mathrm{V}(\mathrm{D}+\propto \mathrm{I})^{-1} \mathrm{~V}^{\prime}$. Computation of $\mathrm{V}(\mathrm{D}+\infty \mathrm{I})^{-1} \mathrm{~V}^{\prime}$ consumes less CPU time than direct inversion mainly because ( $\mathrm{D}+\infty \mathrm{I}$ ) is a diagonal matrix. Although this diagonalization procedure reduces computational time compared to the direct inversion approach, it still involves the calculation of $\mathrm{V}(\mathrm{D}+\propto \mathrm{I})^{-1} \mathrm{~V}^{\prime}$ in each itcration.

Lin [4] applied singular value decomposition to the cocfficient matrix of mixed model equations and used orthogonal matrix V to diagonalize Z'MZ. Although diagonalization of $Z^{\prime} M Z$ involves extensive calculations compared with matrix inversion, it needs to be done only once independently of the number of iterations. After diagonalization, obtaining solutions and estimating variance components are all trivial calculations regardless of the number of iterations, whereas direct inversion approach needs to invert the coefficient matrix in each iteration. Thus Lin's technique will undoubtedly result in a substantial reduction in CPU time compared with the direct inversion approach or the approach of Patterson and Thompson [1].

Lin and Smith [5] applied FG algorithm to transform a multitrait into a unitrait mixed model that has equal design matrices for $t$ traits and contains more than one random effect. The class of models was restricted to those in which the covariance matrices for all random effects including the residual can be diagonalized simultaneously.

All previous studies agreed that inversion of ( $\mathrm{Z}^{\prime} \mathrm{MZ}+\propto \mathrm{I}$ ) is a computational demanding in calculating solution of random effects or in computing REML. estimators of variance components. Appropriateness of an algorithm may change depending on the size of the data and computer capacity.

The purpose of this study is to present nearly simultaneous diagonalization algorithms (F-G and MF-G) as proposed by [6-8] and apply them to the cocfficient matrices of mixed models estimate REML variance components.

## Materials and Methods

## Statistical Model

The mixed linear model that has been used in animal breeding is the following: $y=X b+Z u+c$ where
$y$ is an $n^{*} 1$ data vector of a trait.
$X$ is a known, fixed $n^{*} p$ matrix with $r a n k=r \leqslant \min (n, p)$.
$b$ is a fixed unknown vector.
Z is a known incidence $\mathrm{n}^{*} \mathrm{q}$ matrix.
$\mathbf{u}$ is a nonobservable $q^{*} 1$ random vector (say sire).
e is a $\mathrm{n}^{*} 1$ nonobservable random vector.
$\mathrm{E}(\mathrm{u})=\mathrm{E}(\mathrm{e})=\mathrm{O}, \mathrm{V}(\mathrm{u})=\mathrm{Ae}_{\mathrm{u}}^{2}$ if $\mathrm{A}^{-1}$ (inverse of numerator relationship matrix) is used. otherwise $\mathrm{V}(\mathrm{u})=\mathrm{I} \sigma_{\mathrm{u}}^{2} ; \mathrm{V}(\mathrm{c})=\mathrm{I} \sigma_{\mathrm{e}}^{\mathrm{e}}[9, \mathrm{p} .16]$.

The mixed-model equations (MME) of Henderson [9] are:

$$
\left[\begin{array}{ll}
X^{\prime} X & X^{\prime} Z \\
& \\
Z^{\prime} X & Z^{\prime} Z+\infty I
\end{array}\right]\left[\begin{array}{l}
\hat{6} \\
\hat{u}
\end{array}\right]\left[\begin{array}{l}
X^{\prime} Y \\
\\
Z^{\prime} Y
\end{array}\right]
$$

After absorption of the fixed effects, Henderson's mixed model equations will be $\left(\mathrm{Z}^{\prime} \mathrm{MZ}+\infty \mathrm{I}\right) \mathrm{u}=\mathrm{Z}^{\prime} \mathrm{MY}$ where $\mathrm{M}=\mathbf{I}-\mathrm{X}^{\prime}\left(\mathrm{X}^{\prime} \mathrm{X}\right)^{-1} \mathrm{X}$
and $\propto=\frac{\hat{\sigma}_{e}^{2}}{\hat{\sigma}_{u}^{2}}$
Thus $\hat{\mathbf{u}}=\mathrm{C}^{-1} \mathrm{Z}^{\prime} \mathrm{MY}$ and $\mathrm{C}=\left(\mathrm{Z}^{\prime} \mathrm{MZ}+\infty \mathbf{I}\right), \mathrm{C}$ has the order of q sires and is difficult to compute if $q$ is large.

The REML estimates of sire and error variance components were:

$$
\begin{aligned}
& \hat{\sigma}_{\mathrm{e}}^{2}=\left[\mathrm{Y}^{\prime} \mathrm{MY}-\mathrm{u}^{\prime}\left(\mathrm{Z}^{\prime} \mathrm{MY}\right)\right] /[\mathrm{N}-\operatorname{rank}(\mathrm{x})] \\
& \left.\hat{\sigma}_{\mathrm{u}}^{2}=\mathrm{u}^{\prime} \mathbf{u}+\hat{\sigma}_{\mathrm{e}}^{2} \operatorname{tr}\left(\mathrm{Z}^{\prime} \mathbf{M Z}+\infty \mathbf{I}\right)^{-1}\right] / \mathrm{q},
\end{aligned}
$$

where N is the number of observations and q is the number of sires.
F-G or MF-G algorithms can be applied to diagonalize simultaneously coefficient matrices of mixed model equations. The simplified procedure in calculating REML variance components can be summarized as follows:

1. Accumulate the coefficient matrices $\mathrm{C}_{1}, \ldots, \mathrm{C}_{k}$ and $\mathrm{P}_{1}, \ldots \mathrm{P}_{\mathrm{k}}$, where cach C can be one of the form of $Z^{\prime} M Z,\left(Z^{\prime} M Z+\propto I\right)$ or $\left(Z^{\prime} M Z+\propto A^{-1}\right)$ and each $P_{i}$ is in the form $\mathrm{P}_{\mathrm{i}}=\mathrm{Z}^{\prime} \mathrm{MY}$
2. Apply F-G or MF-G algorithms on each $\mathrm{C}_{\mathrm{i}}$, each with dimension $\mathrm{q}^{*} \mathrm{q}$ to obtain the orthogonal matrix $B$.
3. Compute $\mathrm{B}^{\prime} \mathrm{C}_{\mathrm{i}} \mathrm{B}$ and $\mathrm{B}^{\prime} \mathrm{Z}^{\prime} \mathrm{MY}$.
4. One can apply Gaussian elimination to get an exact solution or create a diagonal matrix $D_{i}=\operatorname{Diag}\left(B^{\prime} C_{i} B\right)$ to compute approximate solution.
5. Examine closeness to diagonality by
a) comparing the diagonal elements and the eigenvalues of the transformed matrix.
b) Computing $\mathrm{Q}(\mathrm{B})=\operatorname{det}\left\{\operatorname{Diag}\left(\mathrm{B}^{\prime} \mathrm{C}_{\mathrm{i}} \mathrm{B}\right)\right\} / \operatorname{det}\left(\mathrm{C}_{\mathrm{i}}\right)$
6. Solve for $u^{*}=\left(D_{1}+\propto I\right)^{-1} B^{\prime} Z^{\prime} M Y$.
7. Compute $\operatorname{tr}\left(Z^{\prime} \mathbf{M A}+\infty\right)^{-1}, \mathbf{u}^{*} \mathbf{u}^{*}$ and $\mathbf{u}^{*} \mathrm{Z}^{\prime} \mathbf{M Y}$.

## F-G algorithm

Flury and Gautschi [6] found that for given $k>1$, positive definite $\mathrm{p}^{*} \mathrm{p}$ matrices $\mathrm{C}_{1} \ldots \ldots, \mathrm{C}_{k}$ and k positive integers $\mathrm{n} 1, \ldots . . \mathrm{nk}$. the algorithm finds an orthogonal matrix $\mathbf{B}$ such that:

$$
\begin{equation*}
\Phi(\mathrm{B}) \frac{\mathrm{k}}{\|_{\mathrm{i}=1}}\left[\operatorname{det}\left(\operatorname{diag}\left(\mathrm{~B}^{\prime} \mathrm{C} \mathrm{~B}\right) / \operatorname{det}\left(\mathrm{C}_{\mathrm{i}}\right)\right]_{\mathrm{i}}{ }_{\mathrm{i}}\right. \text { is minimum } \tag{1}
\end{equation*}
$$

The matrix B brings matrices $\mathrm{C}_{1} \ldots \ldots, \mathrm{C}_{k}$ simultaneously as close to diagonality as possible. Flury [10] showed by using the maximum likelihood estimation of common principal axis in k normal populations that
$\Phi(\mathrm{B})$ is minimum if the following system of equations holds:
where

$$
\begin{equation*}
\operatorname{Eih}_{i h}=b_{h} C_{1} b_{n}(i=1, \ldots \ldots, k ; h=1, \ldots \ldots, p) \tag{3}
\end{equation*}
$$

The F-G algorithm consists of two subroutines, called $F$ and $G$ respectively, which minimize $\Phi(\mathrm{B})$ by iteration on two levels: on the outer level ( F -level) every pair $\left(b_{1}, b_{j}\right)$ of column vectors of the current approximation $B$ to the solution $B$ is rotated, such that equation (3) is satisfied. One iteration step of the $F$-algorithm consists of rotation of all $p(p-1) / 2$ pairs of vectors of $B$. On the inner level (G-level), an orthogonal, $2^{*} 2$ matrix, $Q$ which solves a two dimensional analog of (3), is found by iteration. This matrix defines rotation of a pair of vectors currently being used on the F level, Flury and Gautschi [6, p.171,172].

Clarkson [8] modified the F level of F-G algorithm and improved its performance by reducing the number of operations required for each pair of orthogonal column vectors $B_{p}=\left(b_{j}, b_{1}\right)$ in $B$. An orthogonal matrix $P$ is found such that:

$$
P=\left[\begin{array}{cc}
c & -s \\
& \\
s & \mathrm{c}
\end{array}\right]
$$

where $c$ and $s$ are the $\sin$ and cosin of the rotation angle $\left(c^{2}+s^{2}=1\right)$. Given $c$, the updated versions of vectors $b_{j}, b_{1}$ are computed as $B^{n}=B_{p} P$, that is $b_{j}^{n}=c b_{j}+s b_{1}$ and $b_{1}^{n}=-s b_{j}+c b_{1}$ updated vectors.

In Flurry and Gautschi [6] algorithm, maximum likelihood estimates for $C$ werc found via the " $G$ " step by use of $k$ matrices $T i$, where $\left.T_{i}=b_{j}, b_{1}\right) C_{i}{ }^{\prime} b_{j} b_{1}$ ). Roughly $2 k p^{2}$ operations are required to obtain all $k$ matrices $T_{i}$ for one $F$ step. Since each $F$ iteration must consider all $p(p-1) / 2$ possible pairs of vectors $\left(b_{1}, b_{j}\right)$, the order of $k p^{4}$ operations are required in each $F$ iteration in computing the $T_{i}$ 's. This is the maximum number of operations required by any phase of the F-G algorithm. In MF-

G algorithm, the k multiplication is not utilized in computing $\mathrm{T}_{\mathrm{i}}$, , resulting in a significant increase in performance of the algorithm, Clarkson [8,p.148-149] F-G algorithm, $\mathrm{KP}^{3}$ operations are required per F iteration to update the matrices. $\mathrm{T}_{\mathrm{i}}{ }^{\prime} \mathrm{s}$.

## Numerical Example

Table 1. Example data were adapted from Schaeffer [11].

| Herd-year-Season | Sire eartage | No. of progeny | Total yield/100 |
| :--- | :---: | :---: | :---: |
| 1 | 1 | 4 | 531 |
| 1 | 2 | 3 | 449 |
| 1 | 3 | 3 | 416 |
| 2 | 2 | 3 | 411 |
| 2 | 4 | 2 | 298 |
| 3 | 3 | 4 | 624 |
| 3 | 5 | 6 | 983 |
| 4 | 1 | 2 | 302 |
| 4 | 6 | 3 | 526 |
| 5 | 1 | 2 | 321 |
| 6 | 1 | 2 | 254 |
| 6 | 4 | 4 | 746 |
| 6 | 6 | 2 | 363 |

The model used for analyzing the data contains the fixed effect of herd-year-season and random effect of sire. After absorbing the fixed effect and assuming $\propto=15$ the coefficient matrix is

1) With unrelated sires (say base population):

$$
\begin{array}{rl}
\mathrm{Z}^{\prime} \mathrm{MZ}+\propto \mathbf{I}= & {\left[\begin{array}{rrrrrr}
5.1 & -1.2 & -1.2 & -1.0 & 00.0 & -1.7 \\
-1.2 & 3.3 & -.9 & -1.2 & 00.0 & 00.0 \\
-1.2 & -.9 & 4.5 & 00.0 & -2.4 & 00.0 \\
-1.0 & -1.2 & 00.0 & 3.2 & 00.0 & -1.0 \\
00.0 & 00.0 & -2.4 & 00.0 & 2.4 & 00.0 \\
-1.7 & 00.0 & 00.0 & -1.0 & 00.0 & 02.7
\end{array}\right]+15^{*} \mathrm{I}_{6}} \\
& \left(\mathrm{Z}^{\prime} \mathrm{MY}\right)^{\prime}=(-143.35 \\
15.80-21.6078 .90 & 18.80 \\
51.45
\end{array}
$$

2) With related sires (say first generation):

If the relationship matrix among sires, A is

$$
\mathrm{A}=\left[\begin{array}{llllll}
1.000 & 0.000 & 0.000 & 0.000 & 0.000 & 0.000 \\
0.000 & 1.000 & 0.500 & 0.750 & 0.750 & 0.750 \\
0.000 & 0.500 & 1.000 & 0.750 & 0.750 & 0.750 \\
0.000 & 0.750 & 0.750 & 1.250 & 0.750 & 1.000 \\
0.000 & 0.750 & 0.750 & 0.750 & 1.250 & 1.000 \\
0.000 & 0.750 & 0.750 & 1.000 & 1.000 & 1.375
\end{array}\right]
$$

$\mathrm{Z}_{1}{ }^{\prime} \mathrm{MZ}_{1}=\mathrm{L}^{\prime} \mathrm{Z}^{\prime} \mathrm{MZL}$ and L is a lower triangular matrix such that $\mathrm{A}=\mathrm{L}^{\prime} \mathrm{L}$

$$
\mathrm{Z}_{1}{ }^{\prime} \mathrm{MZ}_{1}+=\mathrm{I}=\left[\begin{array}{rrrrrr}
5.000 & -3.825 & -2.208 & -1.308 & -.601 & -1.041 \\
-3.825 & 3.469 & .617 & .769 & .875 & .781 \\
-2.208 & .617 & 2.756 & .934 & -.475 & .451 \\
-1.308 & .769 & .934 & 1.438 & .088 & .152 \\
-.601 & .875 & -.475 & .088 & 1.538 & .585 \\
-1.041 & .781 & .451 & .152 & .585 & 1.013
\end{array}\right]+15{ }^{*} \mathrm{I}_{6}
$$

F-G and MF-G algorithms were applied on different sets of coefficient matrices:

1) $Z^{\prime} M Z$ and $Z_{1}^{\prime} M Z_{1} \quad$ where $Z_{1}^{\prime}=I^{\prime} Z^{\prime}$.
2) $\left(Z^{\prime} M Z+\infty I\right)$ and $\left(Z_{i}^{\prime} M Z_{i}+\infty I\right)$.
3) $\left(\mathrm{Z}^{\prime} \mathrm{MZ}+\infty \mathrm{I}\right)$ and $\left(\mathrm{Z}_{1}^{\prime} M Z_{1}+\infty \mathrm{A}-1\right)$.

These three sets were chosen as an example to demonstrate simultaneous diagonalization of two coefficient mixed model matrices. Moreover, each set will differ from the other in the magnitude of the diagonal and off-diagonal elements.

Different values of $x=15,50$ and 500 were used. An initial matrix $B=I$ and equal and unequal weights were used in F-G and MF-G to compute an orthogonal matrix $B$ which diagonalizes each set. The matrix $B$ which achieves near diagonality for $x=$ 15 for each set is:

$$
\begin{aligned}
& \mathrm{B}_{1}=\left[\begin{array}{rrrrrr}
.7063 & .6302 & .2524 & .1520 & .0417 & .1247 \\
-.5732 & .7765 & -.2048 & -.1234 & -.0338 & -.1012 \\
.3126 & .0000 & .9457 & -.0673 & -.0184 & -.0552 \\
-.2035 & .0000 & .0000 & -.9783 & -.0120 & -.0359 \\
-.0571 & .0000 & .0000 & .0000 & .9983 & -.0101 \\
-.1739 & .0000 & .0000 & .0000 & .0000 & .9848
\end{array}\right] \\
& \mathrm{B}_{2}=\left[\begin{array}{lrrrrr}
.7756 & .5406 & -.0182 & -.1697 & -.0865 & -.2636 \\
-.5135 & .3068 & -.4031 & -.3400 & -.1857 & -.5740 \\
-.2694 & .3601 & .753 & -.2439 & .4057 & -.0821 \\
-.1743 & .3882 & .1891 & .8000 & -.3562 & -.1280 \\
-.0596 & .3614 & -.4764 & .2772 & .7274 & .1815 \\
-.1679 & .4502 & -.0863 & -.2816 & -.3708 & .7382 \\
& & & & &
\end{array}\right] \\
& \mathrm{B}_{3}=\left[\begin{array}{lrrrrr} 
& & & & -.0055 \\
\hline .9861 & .1497 & .0089 & .0257 & .0666 & -.0055 \\
-.0618 & .3878 & -.3696 & -.6244 & .2930 & -.4832 \\
-.0024 & .4567 & .5241 & -.3064 & -.6492 & -.0291 \\
-.0554 & .3872 & -.3443 & .6610 & -.2953 & -.4521 \\
-.1262 & .4701 & .5331 & .2798 & -.6329 & -.0078 \\
-.0657 & .4977 & -.4313 & -.0185 & -.0206 & -.7491
\end{array}\right]
\end{aligned}
$$

Flury [10] and Flury and Gautschi [6] found the eigenvectors of the diagonalizable matrices, but B can be considered as "compromises" between the eigenvectors of the untransformed matrices.

Tables 2 and 3 show that the trace, $u^{\prime} u$ and $u^{\prime} Z^{\prime} M Y$ of the three sets of transformed matrices are approximate to those of exact solution (direct inversion). Diagonalization of ( $\left.\mathrm{Z}^{\prime} \mathrm{MZ}+\infty \mathrm{I}\right)$ and $\left(\mathrm{Z}_{\mathrm{I}} \mathrm{MZ}_{1}+\propto \mathrm{I}\right)$ gave the nearest results to the exact solution. Dropping the off-diagonal elements from the transformed matrices gave approximate variance components of REML (Tables 2 and 3). One can get the exact solution by inverting the complete transformed matrices which is computationally demanding. The difference between the approximate solution of REML and the exact solution could be narrowed by magnifying the diagonal clements. One can take

Table 2. Approximate estimates of the trace, $u^{\prime} u$ and $u^{\prime} Z^{\prime} M Y$ for different ratio $(\mathbf{R}=\propto \mathbf{I})$ and different sets of matrices.

| $\propto$ | Set | Trace | $\mathbf{u}^{\prime} \mathbf{u}$ | $\mathbf{u}^{\prime} \mathbf{Z}^{\prime} \mathbf{M Y}$ |
| :---: | :---: | :---: | :---: | :---: |
| 15 | DI | . 32979 | 72.166 | 1480.590 |
|  | First | . 32499 | 80.036 | 1558.700 |
|  | Second | . 32899 | 76.763 | 1525.140 |
|  | Third | . 32782 | 80.533 | 1563.790 |
| 50 | DI | . 11232 | 9.854 | 548.000 |
|  | First | . 11213 | 10.227 | 558.282 |
|  | Sccond | . 11229 | 10.070 | 553.923 |
|  | Third | $.11222$ | $10.385$ | 562.415 |
| 500 | DI | . 1192 | . 11925 | 60.299 |
|  | First | .0!192 | . 11974 | 60.456 |
|  | Second | . 01192 | . 11953 | 60.367 |
|  | Third | . 01192 | . 12001 | 48.903 |

DI $=$ direct inversion, First Set $=\left(Z^{\prime} M Z\right),\left(Z^{\prime} M Z\right),\left(Z^{\prime}, M_{1}\right)$, Second Set $=\left(Z^{\prime} M Z+\infty \mathbf{1}\right)$, $\left(Z^{\prime} 1 \mathrm{MZ}_{1}+\infty 1\right)$. Third set $=\left(Z^{\prime} M Z+\infty \mathrm{I}\right),\left(Z^{\prime} \mathrm{MZ}+\infty \mathrm{A}^{\cdot 1}\right)$

Table 3. Approximate estimates of the trace, $u$ ' $u$ and $u^{\prime} Z^{\prime} M Y$ for different ratio $\left(R=x^{\prime} A^{-1}\right)$ and different sets of matrices.

| $x$ | Set | Trace | $u^{\prime} \mathbf{u}$ | $\mathbf{u}^{\prime} \mathbf{Z}^{\prime} \mathbf{M Y}$ |
| :---: | :---: | :---: | :---: | :---: |
| 15 | DI | . 35166 | 75.392 | 1808.720 |
|  | First | . 35121 | 81.569 | 1540.660 |
|  | Second | . 35121 | 78.075 | 1506.120 |
|  | Third | . 380007 | 147.574 | 1624.480 |
| 50 | Dl | . 11459 | 12.460 | 738.140 |
|  | First | . 11455 | 9.999 | 550.488 |
|  | Second | . 11457 | 9.838 | 546.014 |
|  | Third | $.12845$ | 32.043 | 737.552 |
| 500 | DI | . 01194 | . 16873 | 85.931 |
|  | First | $.01194$ | . 11925 | 60.301 |
|  | Second | . 01194 | . 11903 | 60.240 |
|  | rhird | . 01366 | . 43684 | 85.994 |

$D I=$ direct inversion, First Set $=\left(Z^{\prime} M Z\right),\left(Z^{\prime} 1 M Z 1\right)$, Second Sct $=\left(Z^{\prime} M Z+\propto I\right),\left(Z^{\prime} 1 M Z 1+\propto I\right)$, Third set $=\left(Z^{\prime} M Z+\infty I\right),\left(Z^{\prime} M Z+\infty A^{-1}\right)$
advantage of the large diagonal elements relative to small off-diagonal elements and diagonalize mixed model coefficient matrices after adding to diagonal element. Flury and Gautschi [6] found that iterative F-G algorithm converges faster with large diagonal elements. Moreover, Schaeffer [11] found that the iterative solution of large mixed model converges faster with large diagonal elements, and the larger are the diagonals compared to off-diagonal elements in the equations, the faster will be the rate of convergence.

Two criteria must be met to achieve complete diagonality and consequently finding an exact solution:

1) the diagonal elements of the transformed matrices are identical with their respective eigenvalues. Comparison of diagonal elements and their corresponding eigenvalues in the numerical example was given in Table 4 and 5. Diagonal elements and the eigenvalues became close to each other as increased to 500. The agreement of diagonal elements with the corresponding eigenvalues needs to be checked using likelihood ratio test as suggested by [10]. If the diagonal elements and their corresponding cigenvalues differ significantly then approximate estimation should be considered cautiously.
2) $Q(B)=1$ where

$$
Q(B)=\frac{\left|\operatorname{Diag}\left(B^{\prime} C_{i} B\right)\right|}{\left|C_{i}\right|}
$$

$\left|\left(\operatorname{diag}\left(\mathrm{B}^{\prime} \mathrm{C}_{\mathrm{i}} \mathrm{B}\right)\right)\right|=$ the product of all diagonal elements of the matrix inside the parenthesis.
$Q(B)>1$ if the off-diagonal elements deviate from zero. Table 6 shows estimates of $\mathrm{Q}(\mathrm{B})$ for different sets of simultaneous transformation. As $x$ increased to 500 . $Q(B)$ become close to 1 . Diagonalizing $Z^{\prime} M Z$ and $Z_{1}^{\prime} \mathrm{MZ}_{1}$ gave large values of $\mathrm{Q}(\mathrm{B})$, and this is mainly due to the very small determinant of both matrices, $\operatorname{det}\left(Z^{\prime} \mathrm{MZ}\right)=2.5941 \mathrm{E}-13$, and $\operatorname{det}\left(\mathrm{Z}^{\prime}{ }_{\mathrm{I}} \mathrm{MZ}_{1}\right)=-2.1110 \mathrm{E}-15$.

Flury and Gautschi [6] showed that two minima of equation (1) are expected if the matrix has small determinant, i.e. if the essentricity $\left(\frac{\mu_{\max }}{\mu_{\min }}\right.$, where $u$ is the eigenvalue) is high.

One can easily find that sums of squares of the off-diagonals for coefficient matrices are less after applying F-G and mf-G algorithms. Moreover, in terms of abso-

Table 4. Diagonal elements (DE) and eigenvalues (EV) for different sets of coefficient matrices and for different ratio

| Set ${ }^{-1}$ | $\mathrm{R}=15 \mathrm{I}$ |  | $\mathrm{R}=50 \mathrm{I}$ |  | $\mathrm{R}=500 \mathrm{I}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | DE | EV | DE | EV | DE | EV |
| First | 20.738 | 22.016 | 55.738 | 57.016 | 505.738 | 507.016 |
|  | 17.841 | 17.916 | 52.841 | 52.916 | 502.841 | 502.916 |
|  | 19.388 | 20.844 | 54.388 | 55.855 | 504.388 | 505.843 |
|  | 18.298 | 17.916 | 53.298 | 54.358 | 503.298 | 504.358 |
|  | 17.499 | 16.067 | 52.499 | 52.370 | 502.499 | 501.067 |
|  | 17.437 | 15.000 | 52.437 | 50.000 | 502.437 | 500.000 |
| Second | 21.017 | 22.016 | 56.002 | 57.012 | 505.973 | 507.052 |
|  | 15.157 | 15.157 | 50.166 | 49.998 | 500.156 | 499.973 |
|  | 21.269 | 20.842 | 56.267 | 55.843 | 605.325 | 505.842 |
|  | 19.393 | 19.358 | 54.389 | 54.355 | 504.385 | 504.349 |
|  | 16.112 | 16.067 | 51.110 | 51.063 | 501.091 | 501.039 |
|  | 18.253 | 18.253 | 53.253 | 52.915 | 503.290 | 502.039 |
| Third | 20.489 | 21.666 | 55.302 | 57.014 | 505.141 | 507.007 |
|  | 15.460 | 15.062 | 55.842 | 55.846 | 505.824 | 505.849 |
|  | 16.674 | 16.008 | 54.652 | 54.362 | 504.643 | 504.369 |
|  | 19.406 | 19.406 | 50.816 | 49.999 | 501.058 | 500.025 |
|  | 21.054 | 21.100 | 51.723 | 51.067 | 501.744 | 501.065 |
|  | 18.074 | 17.914 | 52.866 | 52.915 | 502.821 | 502.919 |

First Set $=\left(Z^{\prime} M Z\right),\left(Z^{\prime} M_{1} Z_{1}\right)$. Second Set $=\left(Z^{\prime} M Z+\propto I\right),\left(Z^{\prime} 1+\infty I\right),\left(Z^{\prime}{ }^{\prime} M_{1} Z_{1}+\infty I\right)$ Third Set $=$ $\left(Z^{\prime} M Z+\infty I\right),\left(Z^{\prime} M Z+\infty A^{-1}\right)$.

Table 5. Diagonal elements (DE) and eigenvalues (EV) for different sets of coefficient matrices and for different ratio

| Set | $\mathrm{R}=15 \mathrm{~A}^{-1}$ |  | $\mathrm{R}=50 \mathrm{~A}^{-1}$ |  | $R=500 \mathrm{~A}^{-1}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | DE | EV | DE | EV | DE | EV |
| Firs1 | 24.590 | 24.625 | 59.599 | 59.625 | 509.590 | 509.625 |
|  | 15.374 | 15.000 | 50.374 | 50.002 | 500.374 | 500.000 |
|  | 17.038 | 18.014 | 52.038 | 53.014 | 502.038 | 503.014 |
|  | 16.061 | 16.009 | 51.061 | 51.009 | 501.061 | 501.009 |
|  | 16.470 | 16.251 | 51.470 | 51.415 | 501.470 | 501.251 |
|  | 15.783 | 15.415 | 50.783 | 50.415 | 500.783 | 500.415 |

Table 5. Diagonal elements (DE) and eigenvalues (EV) for different sets of coefficient matrices and for different ratio

| Set | $\mathrm{R}=15 \mathrm{~A}^{-1}$ |  | $\mathrm{R}=50 \mathrm{~A}^{-1}$ |  | $R=500 A^{-1}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | DE | EV | DE | EV | DE | EV |
| Second | 24.542 | 24.623 | 59.552 | 59.622 | 509.538 | 509.600 |
|  | 15.660 | 15.416 | 50.645 | 49.996 | 500.622 | 499.600 |
|  | 18.010 | 18.014 | 53.006 | 53.010 | 503.061 | 503.068 |
|  | 15.829 | 16.252 | 50.823 | 51.252 | 500.820 | 501.271 |
|  | 15.581 | 14.998 | 50.581 | 50.414 | 500.562 | 500.383 |
|  | 15.693 | 16.010 | 50.695 | 51.006 | 500.730 | 501.026 |
| Third | 20.360 | 20.318 | 55.202 | 55.202 | 505.130 | 505.130 |
|  | 4.217 | 4.215 | 105.848 | 105.842 | 1005.760 | 1005.760 |
|  | 91.000 | 91.032 | 289.898 | 289.899 | 2856.070 | 2856.070 |
|  | 34.700 | 34.800 | 12.734 | 12.473 | 116.517 | 116.517 |
|  | 32.046 | 31.944 | 101.731 | 101.731 | 1001.740 | 1001.740 |
|  | 43.892 | 43.906 | 134.357 | 134.364 | 1317.420 | 1317.420 |

First Set $=\left(Z^{\prime} M Z\right),\left(Z^{\prime}{ }^{\prime} M Z 1\right)$, Second Set $=\left(Z^{\prime} M Z+I\right),\left(Z_{1}^{\prime}+\infty I\right),\left(Z^{\prime}{ }^{\prime} \mathbf{M Z}_{1}+\infty I\right)$ Third Set $=$ $\left(Z^{\prime} M Z+\infty I\right),\left(Z^{\prime} M Z+\infty A^{-1}\right)$.

Table 6. Estimates of $\mathbf{Q ( B )}$ for different sets of coefficient matrices and for different ratio

| Set | $\mathbf{1 5}$ | $\propto$ | $\mathbf{5 0}$ |
| :--- | :---: | :---: | :---: |
| First | very large | very large | $\mathbf{5 0 0}$ |
| Second | 1.011265 | 1.001335 | verylarge |
| Third | 1.017800 | 1.002850 | 1.000000 |

First Set $=\left(Z^{\prime} M Z\right),\left(Z_{1}^{\prime} M_{1}\right)$, Second Set $=\left(Z^{\prime} M Z+\infty 1\right),\left(Z_{1}^{\prime} M Z_{1}+\infty I\right)$, Third Set $=\left(Z^{\prime} M Z+\infty I\right)$, ( $Z^{\prime} M Z+\infty A^{-1}$ ).
lute value each diagonal entry is larger than the sum of off-diagonal entries in that row i.e.

$$
\left|a_{i j}\right|>\sum_{j=1}^{n}\left|a_{i j}\right| \text { for } j=1,2, \ldots \ldots, n .
$$

At King Saud University IBM computer, the CPU time is combined with output machine time, so it is difficult to define CPU time used by either algorithm for
diagonalizing two ( $6^{*} 6$ ) matrices. However, Flury and Constantine [7] diagonalized two ( $6^{*} 6$ ) matrices on MV 20 computer for MF-G algorithm with CPU time .070 sec onds this compared with .101 seconds required for MF-G algorithm.

## Conclusion

F-G and MF-G algorithms give approximate estimates of variance component of REML. Both algorithms gave the same transformation matrix (B). Equal or unequal weights $\mathrm{n}_{1}$, $\qquad$ , nk can be used to achieve convergence for both algorithms and minimize the deviation from diagonality. Close estimate to the exact solution can be obtained for traits with low heritability (i.e large $\propto$ ) such as reproduction and fitness. Saving in CPU time by using MF-G algorithm becomes more important as the number of sires, animals in animal model, and traits increases.

## F-G algorithm Adopted from Flury and Gautschi [1]

Let $\Phi(B)=\Phi\left(B^{\prime} C_{1} B, \ldots \ldots, B^{\prime} C_{k} B ; n_{1}, \ldots \ldots, n_{k}\right)$, the $F-G$ algorithm yields a converging sequence of orthogonal matrices $B^{o}, B^{1}, \ldots \ldots$ such that $\Phi\left(B^{f+1}\right) \leqslant \Phi\left(B^{f}\right)$. The algorithm proceeds as follow:

## F algorithm

Step Fo Define $B=\left(b_{1}, \ldots \ldots, b_{p}\right) \in 0(P)$ as an initial apporoximation to the orthogonal matrix minimizing $\Phi$, e.g. $\mathrm{B}<---\mathrm{I}$, put $\mathrm{F}<-----0$.

Step $F_{1}:$ Put $B^{(f)}<----B$ and $f<----f+1$
Step $\mathrm{F}_{2}:$ Repeat steps $\mathrm{F}_{21}$ to $\mathrm{F}_{24}$ for all pairs $(1, \mathrm{j}), \mathrm{l} \leqslant 1 \leqslant \mathrm{j}<\mathrm{p}$
Step $\mathrm{F}_{2-1}$ : put $\mathrm{H}_{\mathrm{p}^{+} 2}<----\left(\mathrm{b}_{1}, \mathrm{bj}\right)$ and

$$
\mathrm{T}_{1}<\cdots\left[\begin{array}{cc}
\mathrm{H}_{1}{ }_{\mathrm{L}} \mathrm{C}_{\mathrm{i}} \mathrm{~b}_{1} & \mathrm{~b}^{\prime}{ }_{\mathrm{L}} \mathrm{C}_{\mathrm{i}} \mathrm{~b}_{\mathrm{j}} \\
& \\
\mathrm{~b}_{\mathrm{j}} \mathrm{C}_{\mathrm{i}} \mathrm{~b}_{1} & \mathrm{~b}^{\prime}{ }_{\mathrm{j}} \mathrm{C}_{\mathrm{i}} \mathrm{~b}_{\mathrm{j}}
\end{array}\right]
$$

The $T_{i}$ are p.d.s. and $(i=1, \ldots \ldots, k)$.
Step $F_{2-\mathrm{i}}:$ Perform the $G$ algorithm on $\left(T_{1}, \ldots \ldots, T_{k}\right)$ to get an

$$
\text { orthogonal matrix } Q=\left[\begin{array}{cc}
\cos \propto & -\sin \propto \\
\sin \propto & \cos \propto
\end{array}\right]
$$

Step $\mathbf{F}_{2-3}$ : Put $\mathrm{H}^{*}\left(\mathrm{p}^{*} 2\right)=\left(\mathrm{b}^{*}{ }_{1}, \mathrm{~b}^{*}{ }_{\mathrm{j}}\right)<-----\mathrm{HQ}$ (This is an orthogonal rotation of the two columns of H by an angle $x$ ).

Step $F_{2-4}$ : In the matrix $B$, replace columns $b_{1}$ and $b_{j}$ by $b_{1}^{*}$ and $b^{*}$, respectively and call the new matrix again $B$.

Step $\mathbf{F}_{3}$ : If for some $\underset{\mathrm{p}}{\epsilon}>0, \Phi\left(\mathrm{~B}^{\mathrm{f-1})}-\Phi<\mathrm{Bf}\right)>\epsilon$ holds stop. Otherwise start the next iteration stop at $F_{1}$.

## G-algorithm

This algorithm solves the equation:

$$
\begin{equation*}
\mathrm{q} 1\left[\sum_{\mathrm{i}=1}^{\mathrm{k}} \frac{k_{\mathrm{i} 1}-k_{\mathrm{i} 2}}{\mathrm{nit}_{\mathrm{i}}-\cdots-\cdots-\cdots} \mathrm{T}_{\mathrm{i} 1}\right] \mathrm{k}_{\mathrm{i} 2}=0 \text {, where } \tag{1}
\end{equation*}
$$

$\mathrm{T}_{1}, \ldots \ldots, \mathrm{~T}_{\mathrm{k}}$ are fixed p.d.s. $2^{*} 2$ matrices, $\mathrm{n}_{\mathrm{i}}>0$ are fixed constants, $k_{\mathrm{ij}}=\mathrm{q}^{\prime} \mathrm{j}$ Tiqj $(\mathrm{i}=$ $1, \ldots \ldots, k), j=1,2)$ and $Q=\left(q 1, q_{2}\right)$ is an orthogonal $2^{*} 2$ matrix. The iteration of the sequence of orthogonal matrices $\mathrm{Q}^{\circ}, \mathrm{Q}^{1}, \ldots \ldots$, converging to a solution of the algorithm proceeds as follows:

Step Go: Define $\mathrm{Q}\left(2^{*} 2\right)$ as an initial approximation to the solution on (1).

$$
\mathrm{Q}<------\mathrm{I}_{2} \quad \text { put } \mathrm{g}<-\cdots-\cdots---0
$$

Step G1: Put $\mathrm{O}(\mathrm{g})<------\mathrm{Q}$ and $\mathrm{g}<-\cdots----\mathrm{g}+1$
Step G2: Compute $k_{\mathrm{ij}}$ using the algorithm

Step G3: Compute normalized eigenvectors of $T$. In $Q=\left(q_{1}, q^{2}\right)$, Put $q 1<-\cdots-{ }^{-}$first eigenvector of $\mathrm{T}, \mathrm{q}_{2}<------$ second eigenvector of T .

Step G4: If $\left\|\mathrm{Q}^{q-1}-\mathrm{Q}\right\|<\epsilon$ stop. Otherwise start the next iteration step.

## MF-G algorithm adopted from Clarkson [2]

1. Compute initial matrices $\mathrm{Q}_{1}=\mathrm{B}^{\prime}{ }_{0} \mathrm{C}_{\mathrm{i}} \mathrm{B}_{0}\left(\mathrm{~B}_{0}\right.$ matrix of initial estimates.
2. For column vectors $\left(b_{1}, b_{j}\right)$ of $B$ take the elements of $T_{1}$ as the corresponding diagonal and off-diagonal elements of $Q_{i}$. In other words $t_{i 11}=q_{i j j}, t_{i 22}=q_{i 11}, t_{i 12}=$ $\mathrm{q}_{\mathrm{ij} \mathrm{i}}$ and $\mathrm{t}_{\mathrm{i} 21}=\mathrm{q}_{\mathrm{ij} 1}$
3. Update each matrix $Q_{i}$ from the values $c$ and $s$ computed during the $G$ step as

$$
\begin{aligned}
& q_{i}^{\mathrm{n}} \mathrm{ijj}=\mathrm{c} 2 \mathrm{t}_{\mathrm{i11}}+2 \mathrm{csti}_{12}+\mathrm{s}^{2} \mathrm{ti}_{22} \\
& \mathrm{q}^{2}{ }_{\mathrm{i} 11}=\mathrm{s}^{2} \mathrm{t}_{\mathrm{i} 11}-2 \mathrm{cst}_{\mathrm{i} 12}+\mathrm{c}^{2 \mathrm{t}_{\mathrm{i} 22}} \\
& q^{n}{ }_{i j 1}=c s\left(t_{i 22}-t_{i 11}\right)+\left(c^{2} s^{2}\right) t_{i 12} \\
& q^{\mathrm{n}}{ }_{\mathrm{ij} 1}=\mathrm{q} \mathrm{ij}_{\mathrm{i}} \mathrm{I} \\
& q^{\mathrm{n}}{ }_{\mathrm{iml}}=\mathrm{sq}_{\mathrm{iml}}+\mathrm{sq}_{\mathrm{imj}} \\
& q^{n}{ }_{i m l}=s q_{i m l}+s q_{i m j} \\
& q^{n_{i m j}}=q_{i m j} \\
& q^{n_{i l m}}=q_{i m 1} \\
& \text { where } m=1, \ldots \ldots, p, m=1, m=j
\end{aligned}
$$

4. Update the vectors $\left(\mathrm{b}_{\mathrm{j}}, \mathrm{b}_{1}\right)$ as discussed above using tangent.
5. Go to step 2 with a new pair $(j, 1)$ of indices for vectors $\left(b_{j}, b_{1}\right)$.

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# تطبيق طريقة حساب ف. ج الوترية على أعلى إمكانية محدودة لتقدير مكونات التباين 

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ملخص البحث. يمكن استخدام طريقة حساب ف. ج للعالم فالوري وجاوتش (19^£) لـدساب المصفوفة ب التي تعطي أقل قيمة مكنة للمعادلة

$$
\begin{aligned}
& \text { ك }
\end{aligned}
$$

$$
\begin{aligned}
& \text { أَا هو مقلوب النصفونة أ (مصفوفة القرابة) }
\end{aligned}
$$

المصفوفة ب تحول المصفوفات ج، ج، ج، ، . . . . . . . . . . . . . . . . . . . . . . . .
مصفوفات قريبة من الوترية وبالتالي يمكن إيكاد معلوب المصفونة ز ز م ر + هـ هـ أ أ' . ولثقليل عدد دورات
 بوساطتها يمكن تحويل النصفونة ج د إلى مصفوفة قريبة من الوترية .

أمكن تطبيق طريقتي المساب ف. ج. ج، م ف. جـ على ثلاث بجموعات من المعادلات المزدوجة والتي

 متساوية أو غير متساوية.

